Assertions, Denials Questions, Answers & the Common Ground

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To better understand the speech acts of *assertion* and *denial*, their relationships to *other* speech acts, and connections between speech acts and logical notions, including the Gentzen's sequent calculus.

My Prompt

I want to revisit some themes (and expand on some of the claims) from my 2005 paper "Multiple Conclusions."

My Focus

The behaviour of two kinds of speech acts: (1) *polar* (yes/no) *questions*, and (2) *justification requests*.

My Plan

Assertion and Denial

Polar Questions

Positions and Rules

Justification Requests

ASSERTION AND DENIAL

Multiple Conclusions

$\boldsymbol{X}\succ\boldsymbol{Y}$

Don't *assert* each member of X and *deny* each member of Y.

Structural Rules

$$X, A \succ A, Y$$
 Id $\frac{X \succ A, Y \quad X, A \succ Y}{X \succ Y}$ Cut

These rules govern assertion and denial *as such*.

Defining Rules for Logical Concepts

$$\frac{X, A, B \succ Y}{\overline{X, A \land B \succ Y}} \land Df \quad \frac{X \succ A, B, Y}{\overline{X \succ A \lor B, Y}} \lor Df \quad \frac{X \succ A, Y}{\overline{X, \neg A \succ Y}} \neg Df \quad \frac{X, A \succ B, Y}{\overline{X \succ A \rightarrow B, Y}} \rightarrow Df$$
$$\frac{X \succ A(n), Y}{\overline{X \succ \forall xA(x), Y}} \forall Df \quad \frac{X, A(n) \succ Y}{\overline{X, \exists xA(x) \succ Y}} \exists Df \quad \frac{X, Fa \succ Fb, Y \quad X, Fb \succ Fa, Y}{\overline{X \succ a = b, Y}} = Df$$

Terms \mathfrak{E} *conditions*: the singular term n (in $\forall / \exists Df$) and the predicate F (in =Df) do not appear below the line in those rules.

These rules can be understood as *definitions* of the concepts they introduce (below the double line).

See (Scott 1974; Došen 1980, 1989; Restall 2019).

Assertions, Denials, Questions, Answers, & the Common Ground

In appealing to norms governing assertion...

... I was wading into a pre-existing literature about assertion. A *very large* literature.

Norms for Assertion

It is fruitful to think of assertion as an act governed by *norms*.

For me: Production Norms

Aim to say what is *true*!

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Only say what you know!

For me: Production Norms

Aim to say what is true!

Only say what you know!

Be prepared to *back it up* when requested!

For you: Consumption Norms

The hearer is entitled to re-assert.

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The hearer is entitled to re-assert.

You can refer back to the asserter to *vouch for* the assertion.

To assert is to bid for the content asserted to be added to the COMMON GROUND, the body of information that we (together) take for granted.

To presuppose something is to take it for granted, or at least to act as if one takes it for granted, as background information as common ground among the participants in the conversation. What is most distinctive about this propositional attitude is that it is a social or public attitude: one presupposes that ϕ only if one presupposes that others presuppose it as well.

- "Common Ground" $L \mathscr{CP}(2002)$

What is the relationship between Assertion and Denial?

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This does not help distinguish *denial* from *retraction*, or from other speech acts.

Let's address this issue...

... by examining polar questions, and their answers, in the light of our background interest in assertion and its norms.

POLAR QUESTIONS

Is it the case that p?

This is a distinct speech act with its own norms.

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It raises an *issue*.

There are two ways to settle the issue

Yes

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Yes No

If I say *yes* and you say *no* to some polar question p?, then we DISAGREE.

That is, we take *different* positions on p.

If I say *yes* and you say *no* to some polar question p?, then we DISAGREE.

That is, we take *different* positions on p. There is no *shared* position incorporating both of our answers.

Other responses, like

Other responses, like

maybe

Other responses, like

maybe · I don't know

Other responses, like

maybe · I don't know · I think so

Other responses, like

maybe · I don't know · I think so

are acceptable responses to p?,

Other responses, like

maybe · I don't know · I think so

are acceptable responses to p?, but they don't answer the question. They don't settle the issue of p.

Settling answers are assertions

A yes or a no to p? counts as an assertion.
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(Either answer is governed by all of the assertion norms we've seen.)

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Presumably ¬p.

However, I prefer to think of a *yes* to p? as ruling p *in*, and a *no* to p? as ruling p *out*.

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(Nothing important *here* hangs on this distinction.)

[X : Y]

a pair of sets of sentences

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- > We have ruled *out* everything in Y, the NEGATIVE COMMON GROUND.

[X : Y]

a pair of sets of sentences

- \succ We have ruled *in* everything in X, the POSITIVE COMMON GROUND.
- > We have ruled *out* everything in Y, the NEGATIVE COMMON GROUND.

Think of this as part of the *conversational scoreboard*. There are also our public *individual* commitments, the questions under discussion, and much more.

ABELARD: Astralabe is in the study. ELOISE: No, he is in the kitchen.

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PARTIAL ANSWER

ABELARD: Astralabe is in the study. ELOISE: No, he is in the kitchen.

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WEAK DENIAL

ABELARD: Is Astralabe in the study? ELOISE: No, he is in the kitchen. STRONG DENIAL

ABELARD: Is Astralabe in the study? ELOISE: *No, he is *either* in the kitchen *or* the study.

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PARTIAL ANSWER

Strong and Weak Denial

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- > To weakly deny p is to block the addition of p to the positive common ground, or to bid for its retraction if it is already there.

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Strong and Weak Denial, and the Common Ground

- Strong *or* weak denials of p are appropriate responses to an assertion of p, because the assertion of p is a bid to add p to the positive common ground.
- ≻ A strong denial of p is one way to settle the question p? to settle it negatively.
- A weak denial of p is not generally an appropriate response to the polar question p?, as the polar question does not place p in the positive common ground, and the question is inappropriate if p is already in the positive common ground, so there is no p to block or retract.

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- ≻ STRONG DENIAL: add to the negative common ground.
- \succ strong assertion: add to the positive common ground.
- \succ WEAK DENIAL: retract (or block) from the positive common ground.
- ≻ WEAK ASSERTION: retract (or block) from the negative common ground. — "Perhaps p."

That's one way to understand the relationship between assertion and denial, and how to distinguish strong denial from other negative speech acts.

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ELOISE: No. The interior angles of *this* triangle add up to 180°. Can you prove the general case?

Eloise blocks from the common ground (weakly denies) a *logical consequence* of the common ground (the axioms of geometry), for the same kind of reason we accept for other weak denials.

This would be impossible if the common ground was simply a set of worlds.

Logical Consequence and Strong or Weak Denial

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If $X \succ Y$ is derivable, then it's out of bounds to *strongly assert* each member of X and *strongly deny* each member of Y.

But this example shows that it *need not* be out of bounds to *strongly assert* each member of X and *weakly deny* each member of Y.

Positions

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If $X \succ Y$ is not derivable then [X : Y] is *available*.

A Word on Cut

 $\frac{X \succ A, Y \quad X, A \succ Y}{X \succ Y} \text{ Cut}$

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(After all, isn't the issue A? thereby *implicitly* settled by [X : Y]?)

POSITIONS AND RULES

Defining Rules

 $\frac{X, A, B \succ Y}{\overline{X, A \land B \succ Y}} \land Df \qquad \frac{X \succ A, B, Y}{\overline{X \succ A \lor B, Y}} \lor Df$



Defining Rules





These are kinds of *definitions*, showing how to treat assertions or denials of the *defined* concept in terms of the assertions or denials of their components.

Derivations

$$\frac{\neg p \succ \neg p}{\succ p, \neg p} \neg Df$$
$$\frac{\neg p, \neg p}{\succ p \lor \neg p} \lor Df$$

Derivations

$$\frac{\neg p \succ \neg p}{\succ p, \neg p} \neg Df \qquad \frac{p \succ p}{p, \neg p \succ} \neg Df \qquad \frac{p, \neg p \succ}{p, \neg p \succ} \land Df$$

Derivations

$$\frac{\neg p \succ \neg p}{\succ p, \neg p} \neg_{Df} \qquad \frac{p \succ p}{p, \neg p \succ} \neg_{Df} \\ \frac{\neg p, \neg p \succ}{\neg p \lor \nabla p} \lor_{Df} \qquad \frac{p, \neg p \succ}{p \land \neg p \succ} \land_{Df}$$

$$\frac{p, q \lor r \succ p \land q, q \lor r}{p, q \lor r \succ p \land q, r, q} \lor_{Df} \frac{p \land q, q \lor r \succ p \land q, r}{q, p, q \lor r \succ p \land q, r} \land_{Df}}{\frac{p, q \lor r \succ p \land q, r}{p, q \lor r \succ p \land q, r}}_{Cut}$$

$$\frac{\frac{p, q \lor r \succ p \land q, r}{p, q \lor r \succ (p \land q) \lor r} \lor_{Df}}{p \land (q \lor r) \succ (p \land q) \lor r} \land_{Df}$$

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→ They don't have the same *shape* as proofs. → (Where is the *conclusion* in $p \lor q \succ p, q$?)

- \succ They don't have the same *shape* as proofs.
- \succ (Where is the *conclusion* in p \lor q \succ p, q?)
- ≻ A endsequent X ≻ A doesn't tell you to *infer* A *from* X
 it merely tells you to not assert all members of X and deny A.

Let's make this problem **sharp**

"Well, now, let's take a little bit of the argument in that First Proposition—just *two* steps, and the conclusion drawn from them. Kindly enter them in your note-book. And in order to refer to them conveniently, let's call them A, B, and Z :=

(A) Things that are equal to the same are equal to each other.

(B) The two sides of this Triangle are things that are equal to the same.

(Z) The two sides of this Triangle are equal to each other.

Réaders of Euclid will grant, I suppose, that Z follows logically from A and B, so that any one who accepts A and B as true, must accept Z as true?"

"Undoubtedly! The youngest child in a High School—as.soon as High Schools are invented, which will not be till some two thousand years later—will grant *that*."

"And if some reader had not yet accepted A and B as true, he might still accept the sequence as a valid one, I suppose?"

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"And if some reader had not yet accepted A and B as true, he might still accept the sequence as a valid one, I suppose?"

The Tortoise never asserts A and A \rightarrow Z while denying Z, but he doesn't accept A and A \rightarrow Z *as a reason* for Z.

JUSTIFICATION REQUESTS

- **ABELARD:** Astralabe is in the kitchen.
 - ELOISE: Really?
- ABELARD: I saw him there five minutes ago.
 - ELOISE: OK.

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- ABELARD: Astralabe is in the kitchen.
 - ELOISE: Really?
- ABELARD: I saw him there five minutes ago.
 - ELOISE: Yes, but he was in the study two minutes ago.

Justification Requests and Norms for Assertion

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If I give you permission to ask *me* to vouch for my assertion you should to be able to call me on it.

That's a JUSTIFICATION REQUEST.

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This reason is another assertion [or denial] which must be *granted*, (added to the common ground) in order for the request to be met.

Granting the given reason is *necessary* but not *sufficient* for satisfying the justification request.

Definitions and Justification Requests

ACHILLES So ... this is an *equilateral* triangle.

- TORTOISE I'm sorry, I don't follow, my heroic friend. I've not heard that word before: what does '*equilateral*' mean?
- ACHILLES Oh, that's easy to explain. '*Equilateral*' means having sides of the same length. An *equilateral* triangle is a triangle with all three sides the same length.
- TORTOISE OK. That sounds good. You may continue with your reasoning.
- ACHILLES Well, as I was saying, the sides of this triangle are all one cubit in length, so it is an equilateral triangle.
- **TORTOISE** Perhaps you will forgive me, Achilles, but I still don't follow. I grant to you that the sides of this triangle all have the same length. I fail to see, however, that it *follows* that it is an equilateral triangle. Could you explain why it is?

Definitions and Justification Requests

If I accept the definition $A =_{df} B$, then I should accept granting A as meeting a justification request for the assertion of B and ruling out A as meeting a justification request for B's denial and *vice versa*.

A failure to accept this is a sign that I have not mastered the definition.

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 $\frac{X, A, B \succ Y}{\overline{X, A \land B, \succ Y}} \land Df$

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$$\frac{X, A, B \succ Y}{\overline{X, A \land B, \succ Y}} \land Df$$

It is a mistake to grant A and grant B and to look for something more to discharge a justification request for an assertion of $A \land B$, if you take $\land Df$ as a *definition*.

$$\frac{X, A \succ B, Y}{\overline{X \succ A \rightarrow B, Y}} \rightarrow Df$$

$$\frac{X, A \succ B, Y}{\overline{X \succ A \rightarrow B, Y}} \rightarrow Df$$

It is a mistake to rule A in and rule B out and to look for something more to discharge a justification request for a denial of $A \rightarrow B$ if you accept $\rightarrow Df$ as a definition. Justification Requests, Defining Rules and Derivations

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Consider this *focussed* derivation:

$$\frac{A \to Z \succ A \to Z}{A \to Z, A \succ Z} \to Df$$

 \succ Read the *premise* as telling us that in a position in which A → Z is already ruled in, we have an answer to the justification request for A → Z's assertion.
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- \succ Read the *premise* as telling us that in a position in which A → Z is already ruled in, we have an answer to the justification request for A → Z's assertion.
- ≻ Then applying $\rightarrow Df$ we see why we have an answer to the request concerning Z's assertion, in a context in which A \rightarrow Z and A have both been ruled in. (In granting A \rightarrow Z and A we have settled Z positively. Its denial is ruled out, since to assert A and deny Z amounts to denying A \rightarrow Z.)

Focussed Derivations and Justification Requests

SLOGAN: A *derivation* of $X \succ A$, Y shows us how to meet a justification request for the assertion of A in any available position extending [X : Y].

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A derivation of X, $A \succ Y$ shows us how to meet a justification request for the denial of A in any available position extending [X : Y].

(Note: it's the *derivation* that shows how to meet the justification request, not the mere validity of the sequent.)

Details

See the handout for details on derivations with *focus* and meeting *justification requests*.

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- \succ A derivation of a sequent X \succ A, Y [X, A \succ Y] can be transformed into a *procedure* for meeting a justification request for an assertion of A [denial of A] in any available position, appealing only what is granted in [X : Y], and to the defining rules used in that derivation.

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- ≻ The *negative* view of the bounds is seen in the clash between assertion and denial, and the *positive* view is found in the answers we can give to justification requests.
- ≻ Adopting *defining rules* is one way to be *very* precise about the norms governing the concepts so defined, combining *safety*, *univocity* and *expressive power*.

THANK YOU!

http://consequently.org/presentation/2021/
assertion-denial-qa-common-ground-bristol