

# DEFINING QUANTIFIERS

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TOPICS IN FREE LOGIC ✧ MCMP ✧ 16 NOVEMBER 2024

1. DEFINING RULES

2. QUANTIFIERS & GENERALITY

3. DEFINING RULES for FREE QUANTIFIERS

4. GOING WIDER, TOO...

5. WHAT IS A QUANTIFIER?

on TERMS & VALUES

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# (HOW) DO SEQUENT RULES CHARACTERISE CONCEPTS?

$$\frac{X \vdash A, \gamma}{X \vdash A \text{ tonk } B, \gamma} \text{ tonkR}$$

$$\frac{X, B \vdash \gamma}{X, A \text{ tonk } B \vdash \gamma} \text{ tonkL}$$

TOO STRONG!

$$\frac{p \vdash p \text{ tonk } q \quad p \text{ tonk } q \vdash q}{p \vdash q}$$

$$\frac{X \vdash A, \gamma \quad X \vdash B, \gamma}{X \vdash A \text{ plink } B, \gamma} \text{ plinkR}$$

$$\frac{X, A \vdash \gamma \quad X, B \vdash \gamma}{X, A \text{ plink } B \vdash \gamma} \text{ plinkL}$$

TOO WEAK!

A plink B could be  $A \wedge B$   
or  $A \vee B$  or anything  
in between!

$$\frac{X, A \vdash B, \gamma}{X \vdash A \rightarrow B, \gamma} \rightarrow R$$

$$\frac{X \vdash A, \gamma \quad X', B \vdash \gamma'}{X, X', A \rightarrow B \vdash \gamma, \gamma'} \rightarrow L$$

JUST RIGHT

# A DEFINITION

$x$  is equilateral iff  $x$  has three sides  
& each side has equal length.

# A DEFINITION

as an invertible rule

$x$  has three sides                      each side of  $x$  has equal length.

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$x$  is equilateral

# DEFINING THE CONDITIONAL

$$\frac{X, A \vdash B, \gamma}{X \vdash A \rightarrow B, \gamma} \rightarrow_{DF}$$

This is clearly **uniquely defining** (if  $\rightarrow_1, \neq \rightarrow_2$  are defined using the same rule, they are interchangeable)

$$\frac{X \vdash A \rightarrow_1 B, \gamma}{X, A \vdash B, \gamma} \rightarrow_{DF \uparrow}$$

$$\frac{X, A \vdash B, \gamma}{X \vdash A \rightarrow_2 B, \gamma} \rightarrow_{DF \downarrow}$$

$$\frac{\frac{}{A \rightarrow_2 B \vdash A \rightarrow_2 B} \text{Id}}{A \rightarrow_2 B, A \vdash B} \rightarrow_{DF \uparrow}}{A \rightarrow_2 B \vdash A \rightarrow_1 B} \rightarrow_{DF \downarrow}$$

$$\frac{A \rightarrow_2 B \vdash A \rightarrow_1 B \quad X, A \rightarrow_1 B \vdash \gamma}{X, A \rightarrow_2 B \vdash \gamma} \text{Cut}$$

# DEFINING THE CONDITIONAL

$$\frac{\frac{X, A \multimap B, Y}{\phantom{X \multimap A \rightarrow B, Y}} \rightarrow \text{DF}}{X \multimap A \rightarrow B, Y}$$

This is also conservatively extending — the usual left/right rules can be recovered from  $\rightarrow \text{DF}$  using Cut & ID.

$$\frac{\frac{\frac{\frac{}{A \rightarrow B \multimap A \rightarrow B} \text{Id}}{\phantom{A \rightarrow B \multimap A \rightarrow B}} \rightarrow \text{DF} \uparrow}{X_1 \multimap A, Y_1 \quad A \rightarrow B, A \multimap B} \text{Cut}_A}{X_1, A \rightarrow B \multimap B, Y_1} \text{Cut}_B}{X_1, X_2, A \rightarrow B \multimap Y_1, Y_2} \text{Cut}_B$$

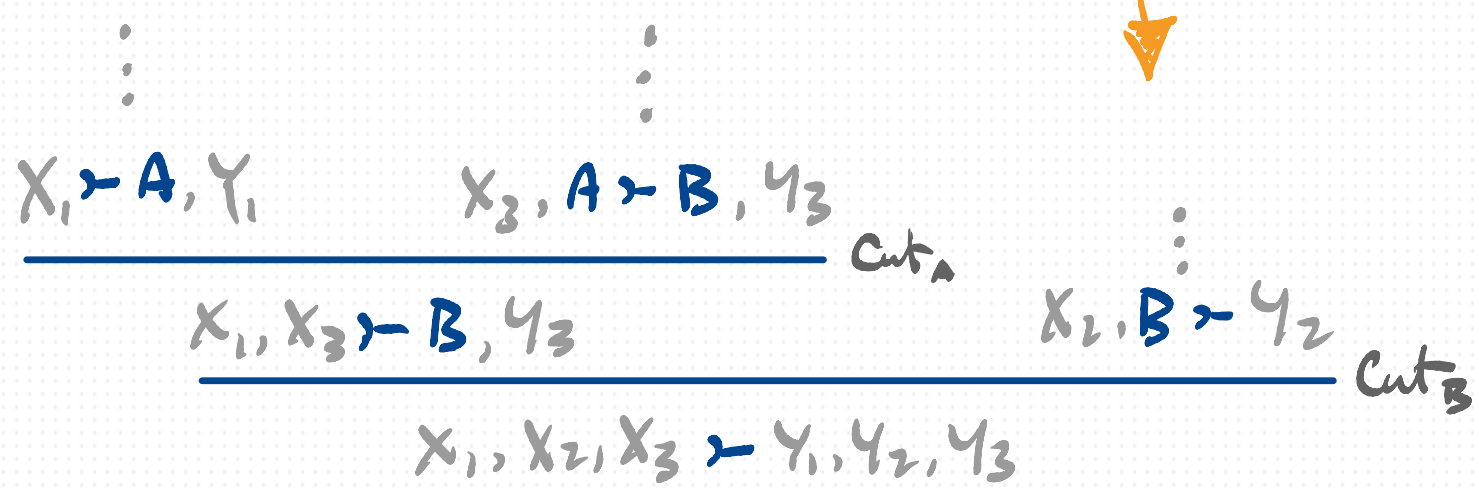
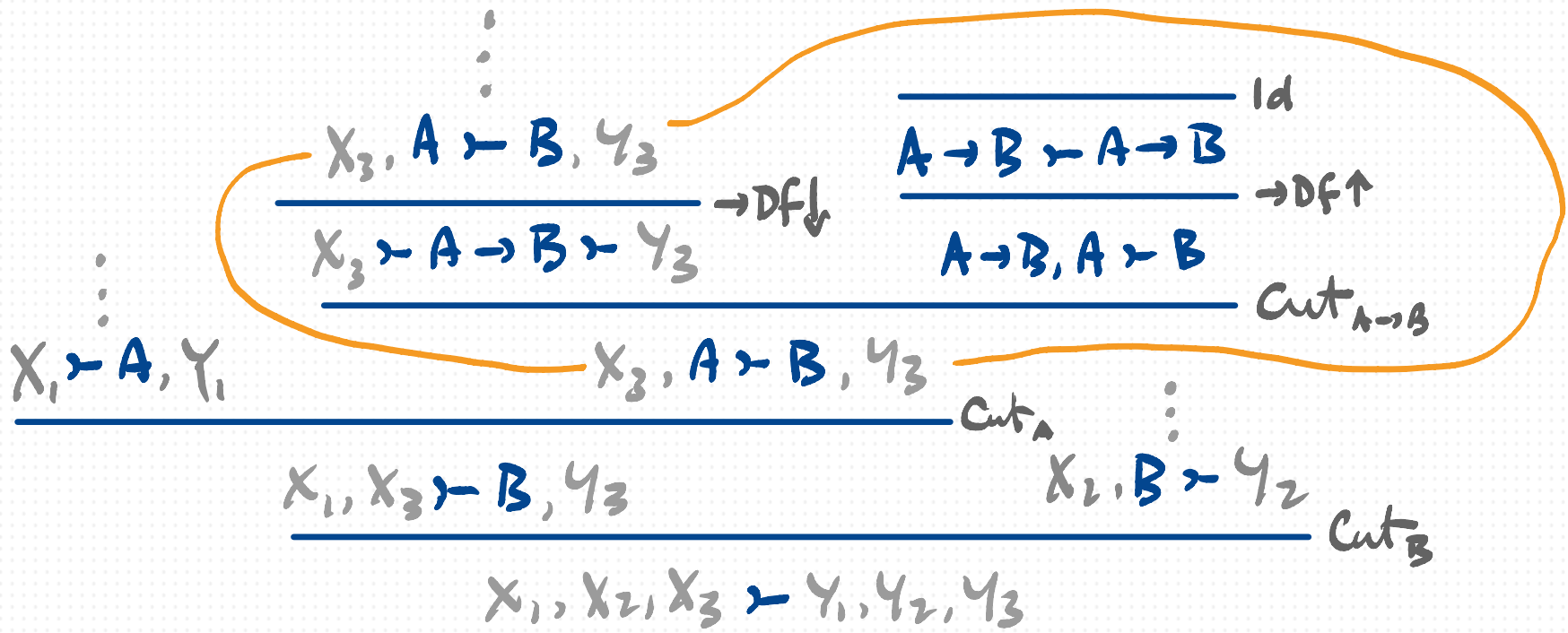


# DEFINING THE CONDITIONAL

$$\frac{\frac{X, A \succ B, Y}{\phantom{X \succ A \rightarrow B, Y}} \rightarrow \text{DF}}{X \succ A \rightarrow B, Y}$$

This is also conservatively extending — the usual left/right rules can be recovered from  $\rightarrow \text{DF}$  using Cut & ID, and the principal cut-reduction step unwinds this definition.

$$\begin{array}{c}
 \vdots \\
 \frac{X_3, A \succ B, Y_3}{\phantom{X_3 \succ A \rightarrow B, Y_3}} \rightarrow \text{DF} \downarrow \\
 \hline
 X_3 \succ A \rightarrow B, Y_3 \\
 \hline
 \vdots \\
 \frac{\frac{\frac{\frac{\frac{A \rightarrow B \succ A \rightarrow B}{\phantom{A \rightarrow B \succ A \rightarrow B}} \text{Id}}{A \rightarrow B \succ A \rightarrow B} \rightarrow \text{DF} \uparrow}}{X_1 \succ A, Y_1 \quad A \rightarrow B, A \succ B} \text{Cut}_A}{\phantom{X_1, X_2, A \rightarrow B \succ Y_1, Y_2}} \text{Cut}_B \\
 \hline
 X_1, X_2, A \rightarrow B \succ Y_1, Y_2 \\
 \hline
 \text{Cut}_{A \rightarrow B} \\
 \hline
 X_1, X_2, X_3 \succ Y_1, Y_2, Y_3
 \end{array}$$



THIS IS **TOTALLY** GENERAL

(it works like this for **any** invertible defining rule, in the presence or absence of contraction & weakening, in different sequent structures.)

$$\frac{X \vdash A, \gamma \quad X \vdash B, \gamma}{X \vdash A \wedge B, \gamma} \wedge Df$$

$$\frac{X, A, B \vdash \gamma}{X, A \otimes B \vdash \gamma} \otimes Df$$

$$\frac{X, A \vdash \gamma \quad X \vdash B, \gamma}{X \vdash A \supset B, \gamma} \supset Df$$

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# DEFINING RULES for QUANTIFIERS

$$\frac{X \supset A(n), Y}{\text{}} \forall \text{df}$$

$$X \supset \forall x A(x), Y$$

$$\frac{X, A(n) \supset Y}{\text{}} \exists \text{df}$$

$$X, \exists x A(x) \supset Y$$

The term  $n$  must be absent in the lower sequent.

$n$  is an **eigenvariable** (more on what this means, soon).

# RECOVERING THE USUAL LEFT/RIGHT RULES

$$\begin{array}{c}
 \text{VR} \\
 \frac{X \succ A(n), Y}{X \succ \forall x A(x), Y} \text{VDF} \downarrow
 \end{array}$$
  

$$\frac{X \succ A(n), Y}{X \succ \forall x A(x), Y} \text{VDF}$$
  

$$\frac{\frac{\frac{\frac{}{\forall x A(x) \succ \forall x A(x)} \text{Id}}{\forall x A(x) \succ A(n)} \text{VDF} \uparrow}}{\forall x A(x) \succ A(t)} \text{Spec}^n}{X, \forall x A(x) \succ Y} \text{Cut}$$
  

$$X, A(t) \succ Y$$

( $t$  can be any term at all, not necessarily an eigenvariable)

# THE SPECULISATION<sup>n</sup><sub>t</sub> RULE

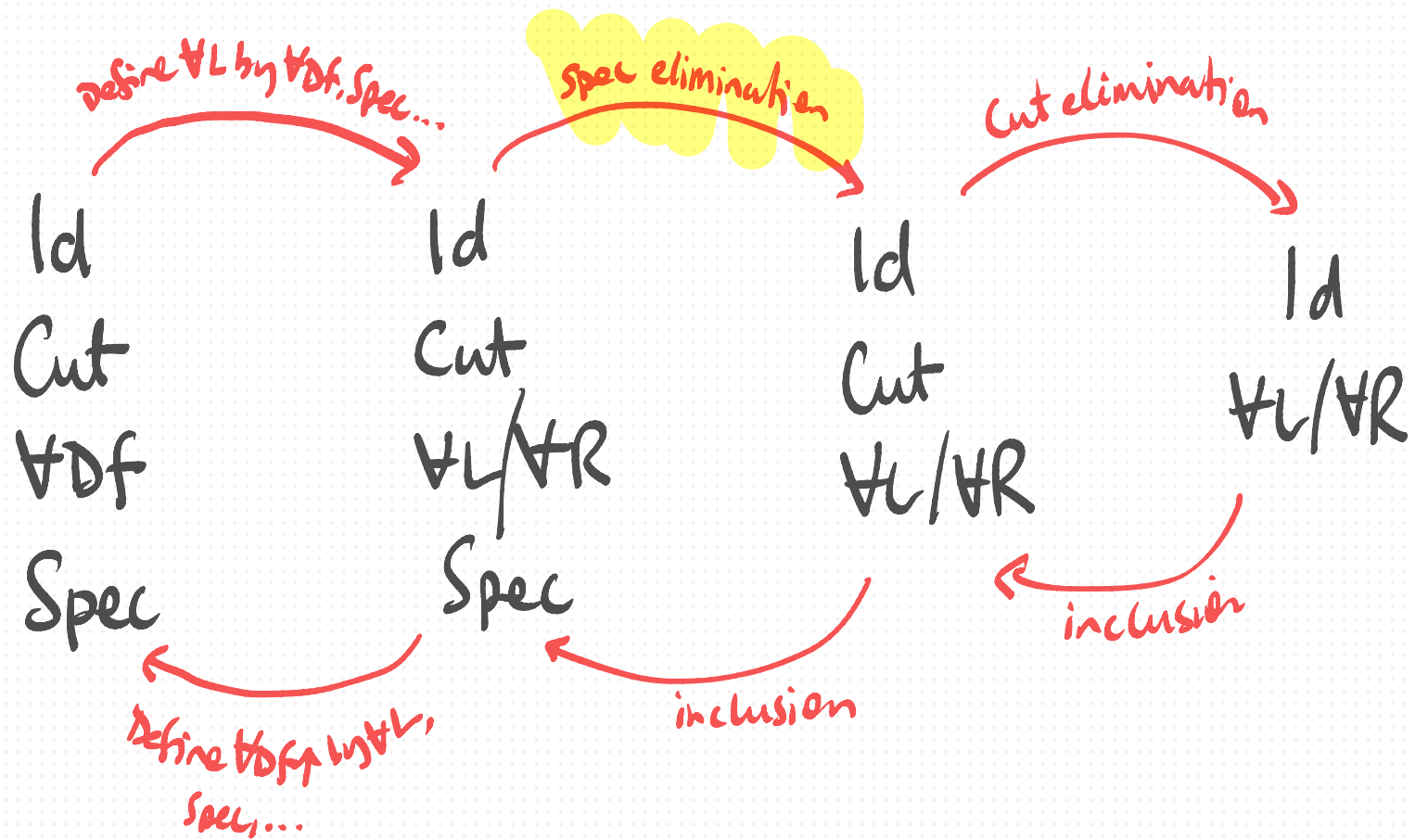
$$\frac{X \supset Y}{[X \supset Y]_t^n} \text{Spec}_t^n$$

If a sequent obtains concerning some eigenvariable  $n$ ,  
it also obtains for the term  $t$ .

Eigenvariables are **inferentially general** among terms.

$\text{Spec}_t^n$  is **admissible** in Gentzen's sequent calculus.

# EQUIVALENT FORMULATIONS



[See my "Generality of Existence 1", RSL (2019) for details.]



# SPEC ELIMINATION

$$\frac{\begin{array}{c} \delta \\ \vdots \\ X \succ Y \end{array}}{[X \succ Y]_t^n} \text{Spec}_t^n \rightsquigarrow \begin{array}{c} \delta_t^n \\ \vdots \\ [X \succ Y]_t^n \end{array}$$

It's easy to see that each of the rules in Gentzen's system are closed under specialisation — ie there is nothing inferentially special in an eigenvariable as conclusion.

( $\forall df \uparrow$  violates this constraint — hence the requirement to impose Spec.)

# DEFINING RULES

- ... specify concepts that are available (conservative) & determinate (unique) over a basic structural context
- ... & for quantifiers, these rules use prior notions of substitution & inferential generality.

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# FREE LOGIC, DEFINED & UNDEFINED TERMS

$n$  is a number.

$n/m$  is not necessarily a number.

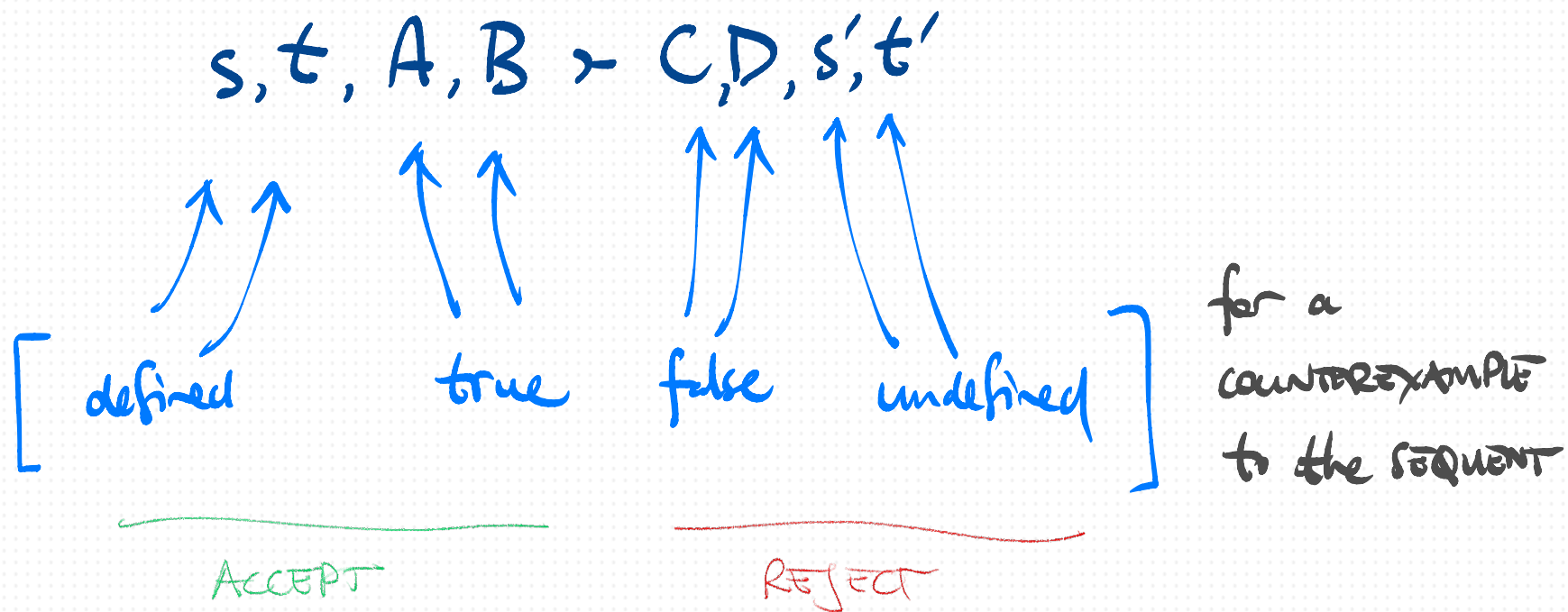
$n/m$  is not defined when  $m=0$ .  
(does not exist)

' $n/m$ ' is a term, whatever values  $n$  &  $m$  take.

$$\forall x (\neg \exists y (y = n/x) \leftrightarrow x=0)$$

# THE BASIC STRUCTURAL CONTEXT

## RULING TERMS IN & RULING THEM OUT



$$\frac{}{t, X \succ Y, t} \text{ t-Id}$$

$$\frac{X \succ Y, t \quad t, X' \succ Y'}{X, X' \succ Y, Y'} \text{ t-Cut}$$

# DEFINING DEFINEDNESS

$$\frac{X \succ Y, t}{X \succ Y, t \downarrow} \downarrow \text{df}$$

& equivalently, ...

$$\frac{X \succ Y, t}{X \succ Y, t \downarrow} \downarrow R$$

$$\frac{X, t \succ Y}{X, t \downarrow \succ Y} \downarrow L$$

# POSSIBLE CONDITIONS ON PREDICATES & FUNCTION SYMBOLS

$$\frac{t_i, X \succ Y}{F t_1 \dots t_n, X \succ Y} \text{ FL}$$

$$\frac{t_i, X \succ Y}{f t_1 \dots t_n, X \succ Y} \text{ FL}$$

(If we grant these, then **nonexistence** is not a predicate in this sense, since we have  $\neg \frac{1}{0} \downarrow$ , while  $\frac{1}{0}$  is not defined.)

# DEFINING RULES FOR $\forall/\exists$ with existential commitment

$$\frac{n, X \supset Y, A(n)}{\frac{}{X \supset Y, \forall x A(x)}} \forall \text{DF}$$

$$\frac{n, A(n), X \supset Y}{\frac{}{\exists x A(x), X \supset Y}} \exists \text{DF}$$

$$\frac{}{\forall x A(x) \supset \forall x A(x)} \text{Id}$$

$$\frac{n, \forall x A(x) \supset A(n)}{\frac{}{t, \forall x A(x) \supset A(t)}} \forall \text{DF} \uparrow$$

$$\frac{t, \forall x A(x) \supset A(t)}{\frac{}{X', A(t) \supset Y'}} \text{Spec}_t^n$$

$$\frac{}{X', A(t) \supset Y'} \text{Cut}$$

$$\frac{X \supset Y, t \quad t, X', \forall x A(x) \supset Y'}{\frac{}{X, X', \forall x A(x) \supset Y, Y'}} t \text{Cut}$$

$$X, X', \forall x A(x) \supset Y, Y'$$



# DEFINING RULES FOR $\forall/\exists$ with existential commitment

$$\frac{n, X \supset Y, A(n)}{\underline{\underline{X \supset Y, \forall x A(x)}}} \forall \text{DF}$$

$$\frac{n, A(n), X \supset Y}{\underline{\underline{\exists x A(x), X \supset Y}}} \exists \text{DF}$$

$$X' \supset A(t), Y'$$

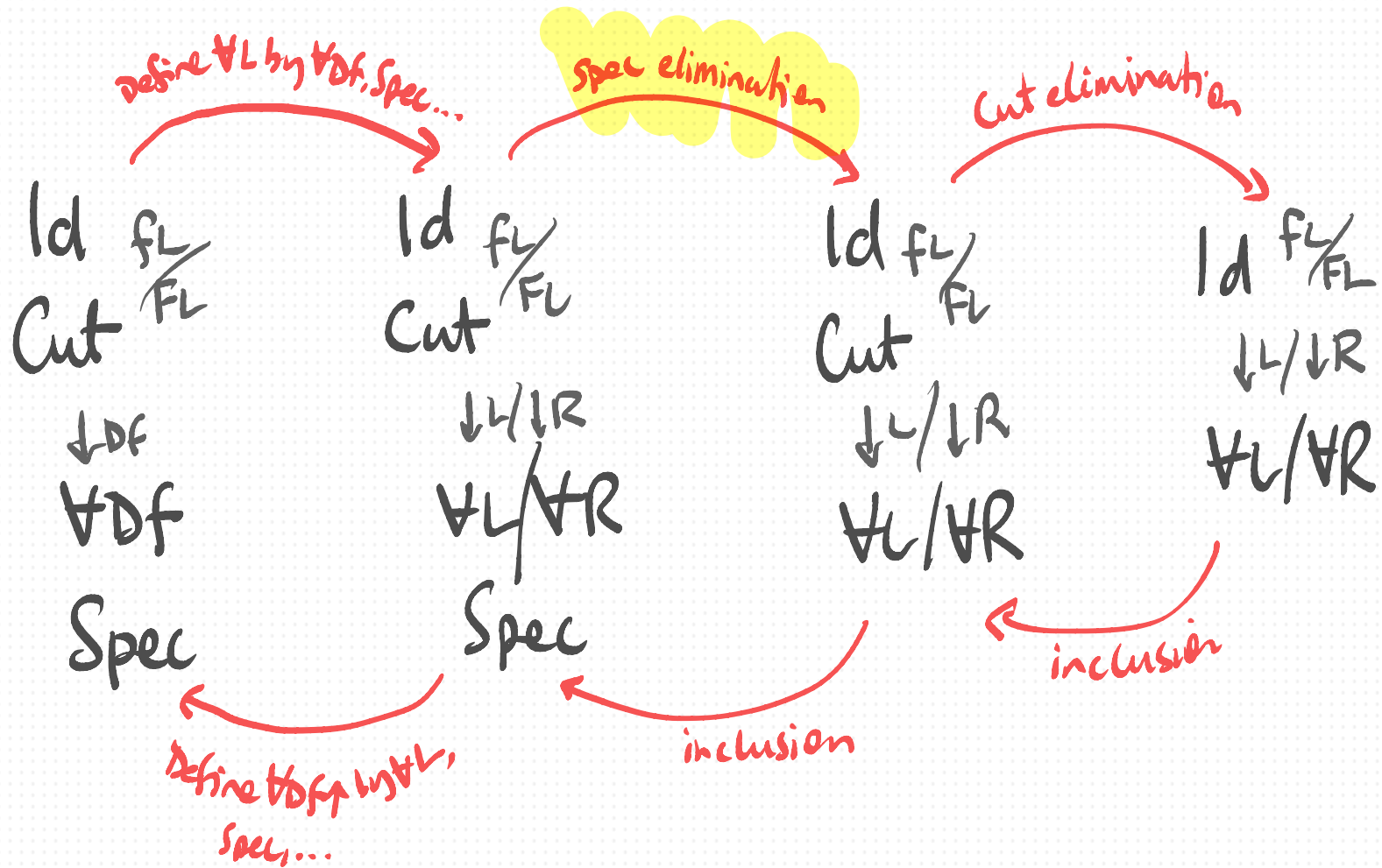
$$\frac{\underline{\underline{\exists x A(x) \supset \exists x A(x)}}}{\underline{\underline{n, A(n) \supset \exists x A(x)}}} \text{Id}$$

$$\frac{\underline{\underline{n, A(n) \supset \exists x A(x)}}}{\underline{\underline{t, A(t) \supset \exists x A(x)}}} \text{Df}\uparrow$$

$$\frac{\underline{\underline{t, A(t) \supset \exists x A(x)}}}{\underline{\underline{t, A(t) \supset \exists x A(x)}}} \text{Spec}_t^n$$

$$\frac{X \supset Y, t \quad \underline{\underline{t, X' \supset \forall x A(x), Y'}}}{\underline{\underline{X, X', \forall x A(x) \supset Y, Y'}}} t \text{Cut}$$

# THE EQUIVALENCES STILL HOLD...



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BUT WE ARE FREE TO DEFINE MORE...

$$\frac{\frac{n, X \supset Y, A(n)}{\text{VDF}}}{X \supset Y, \forall x A(x)}$$

With Existential Commitment

$$\frac{\frac{X \supset Y, A(n)}{\text{PIDF}}}{X \supset Y, \prod x A(x)}$$

without Existential Commitment

$$\frac{\frac{\frac{X \supset Y, \forall x A(x)}{\text{VDF}}}{n, X \supset Y, A(n)}{\text{IDF}}}{\frac{n \downarrow, X \supset Y, A(n)}{\text{ODF}}}{\frac{X \supset Y, n \downarrow \rightarrow A(n)}{\text{PIDF}}}{X \supset Y, \prod x (x \downarrow \rightarrow A(n))}$$

BUT WE ARE FREE TO DEFINE MORE...

$$\frac{n, A(n), X \succ Y}{\exists x A(x), X \succ Y} \exists DF$$

$$\frac{A(n), X \succ Y}{\sum x A(x), X \succ Y} \Sigma DF$$

$$\frac{\sum x \neg x \downarrow \succ \sum x \neg x \downarrow}{\neg n \downarrow \succ \sum x \neg x \downarrow} \Sigma DF \uparrow$$

$$\frac{\neg n \downarrow \succ \sum x \neg x \downarrow}{\neg \frac{1}{0} \downarrow \succ \sum x \neg x \downarrow} \text{Spec}^n \frac{1}{0}$$

Are outer quantifiers just  
as acceptable as inner quantifiers?

Do they involve the same sort of  
{ ontological  
ideological  
theoretical } commitment?

Not necessarily  
There are differences...

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# LOOKING CLOSER...

$$\frac{n, X \succ Y, A(n)}{X \succ Y, \forall x A(x)} \text{VDF}$$

$$\frac{X \succ Y, A(n)}{X \succ Y, \prod x A(x)} \text{PIVDF}$$

this  $n$  is  $\left\{ \begin{array}{l} \text{an eigenvariable} \\ \text{an inferentially general singular term.} \end{array} \right.$

# WHAT IS A VARIABLE?

What is the meaning of ' $x$ ' in  $\forall x A(x)$ ?

"Everything is  $A$ ..."      "For every value  $x$  can take..."

What is the meaning of ' $n$ ' in  $A(n)$ ?

- A variable of some sort?
- A pronoun? a demonstrative expression?
- A singular term of some kind?

Should these be connected?

## IN FAVOUR OF CONNECTION...

$$\frac{x, X \supset Y, A(x)}{\underline{\underline{X \supset Y, \forall x A(x)}}} \forall I$$

$$\frac{x, A(x), X \supset Y}{\underline{\underline{\exists x A(x), X \supset Y}}} \exists I$$

The defining rules become **compositional**, just like the other defining rules.

BUT THEN...

$$\frac{x, X \supset Y, A(x)}{X \supset Y, \forall x A(x)} \forall I$$

$$\frac{x, A(x), X \supset Y}{\exists x A(x), X \supset Y} \exists I$$

THIS SEEMS REDUNDANT!

What would it mean for the variable  $x$   
to be undefined in this context?

Isn't using a variable tantamount to  
treating it as having some value or other?

(E.g. Feferman takes  $x \downarrow$  to be a theorem.)

# ALTERNATE FORMULATION

- Variables are always treated as defined

$$\frac{}{X \succ Y, x} \quad x \downarrow$$

- They are no longer inferentially general but are general among the defined terms.

$$\frac{X \succ Y}{t, [X \succ Y]_t} \quad \text{Spec}_t^x$$

- Variables are used for quantification, not general singular terms.

$$\frac{X \succ Y, A(x)}{X \succ Y, \forall x A(x)} \quad \forall \text{df}_x$$

# EQUIVALENCE

Derivations in the two systems are inter-translatable

Id, Cut

FL fL

Spec<sub>t</sub><sup>n</sup>

Connective Df...

$\forall Df_n \exists Df_n$

$x, y, \dots$  only ever bound.

Eigenvariables  $n$  occur free.

Eigenvariable with existential  
commitment

Id, Cut

FL fL

Spec<sub>t</sub><sup>x</sup>  $x \downarrow$

Connective Df...

$\forall Df_x \exists Df_x$

$x, y, \dots$  occur free

No eigenvariables

Variable



## THE UPSHOT

This formulation provides a context in which inner quantification is definable, using the resources available, but outer quantification (defined in the obvious way) uses extra resources.

This is not to say that outer quantification is meaningless, of course!

# LESSONS & FURTHER QUESTIONS

- ▶ Quantification involves a number of connected issues  
SUBSTITUTION, GENERALITY, DEFINEDNESS, VARIABLES/VALUES, ETC.
- ▶ Choosing a structural context for deduction means taking sides on these issues.
- ▶ Quantification into sentence position? predicate position?  
— do similar issues apply there?
- ▶ Modal calculi?  
— what of singular terms that take values in some worlds and not others?



